



Finanziato  
dall'Unione europea  
NextGenerationEU



Ministero  
dell'Università  
e della Ricerca



Italiadomani  
PIANO NAZIONALE  
DI RIPRESA E RESILIENZA



SERICS  
SECURITY AND RIGHTS IN THE CYBERSPACE



# CrypTO CONFERENCE



Politecnico  
di Torino

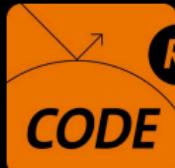
 **Telsy** | A TIM  
ENTERPRISE  
BRAND

 **DISMA**

# REPRESENTING ISOGENIES VIA (CLASSICAL, CANONICAL, ATKIN...) MODULAR POLYNOMIALS

Joint work with Thomas den Hollander, Marc Houben,  
Sören Kleine, Daniel Slamanig, Sebastian A. Spindler

Marzio Mula • May 23, 2025

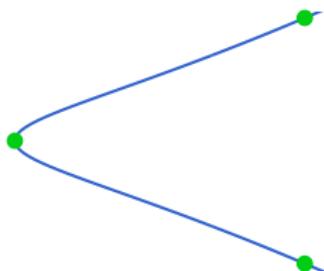


**Research Institute  
Cyber Defence**  
Universität der Bundeswehr München

## IsoGENIES

An **isogeny** is a map between elliptic curves.

$$E : y^2 = x^3 + 13x - 34$$



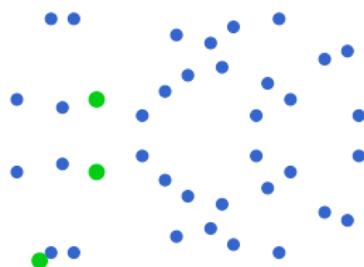
$$E' : y^2 = x^3 - 7x - 6$$



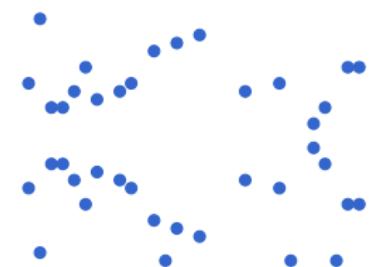
## IsoGENIES

An **isogeny** is a map between elliptic curves.

$$E : y^2 = x^3 + 13x - 34$$



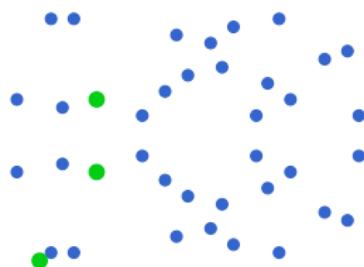
$$E' : y^2 = x^3 - 7x - 6$$



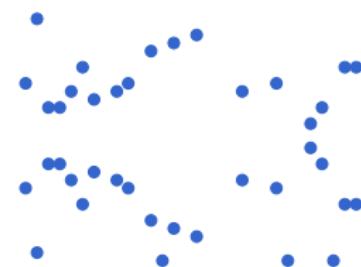
## IsoGENIES

An **isogeny** is a map between elliptic curves.

$$E : y^2 = x^3 + 13x - 34$$



$$E' : y^2 = x^3 - 7x - 6$$



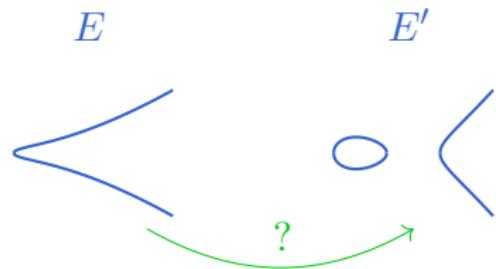
The **degree** of a (separable) isogeny is the cardinality of its kernel.

## A HARD MATHEMATICAL PROBLEM

$p =$  a prime of cryptographic size

$E, E'$  = supersingular elliptic curves over  $\mathbb{F}_{p^2}$

**Problem (ISOGENYPATH):** Given  $E$  and  $E'$ , find an isogeny  $\varphi: E \rightarrow E'$ .

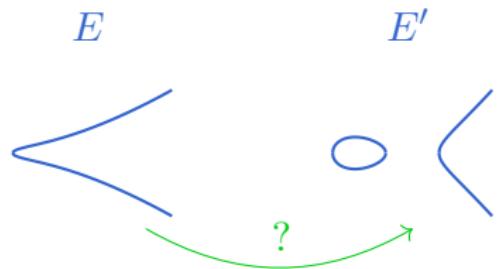


## A HARD MATHEMATICAL PROBLEM

$p =$  a prime of cryptographic size

$E, E'$  = supersingular elliptic curves over  $\mathbb{F}_{p^2}$

**Problem** (`ISOGENYPATH`): Given  $E$  and  $E'$ , find an isogeny  $\varphi: E \rightarrow E'$ .



`ISOGENYPATH` is considered hard even for quantum computers.

## A HARD MATHEMATICAL PROBLEM

$p =$  a prime of cryptographic size

$E, E'$  = supersingular elliptic curves over  $\mathbb{F}_{p^2}$

**Problem** (ISOGENYPATH): Given  $E$  and  $E'$ , find an isogeny  $\varphi: E \rightarrow E'$ .

public parameter      public key



ISOGENYPATH is considered hard even for quantum computers.



Can be used for **post-quantum cryptography!**

- Key exchanges (CSIDH, SCALLOP, SCALLOP HD...)
- Signatures (SQISign, PRISM)

## PROVING KNOWLEDGE OF AN ISOGENY

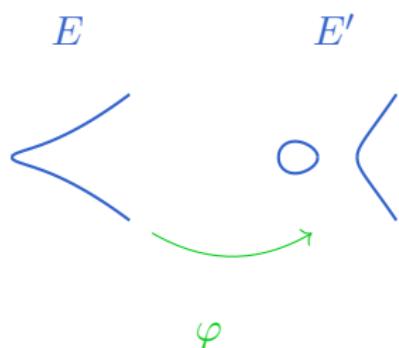
I know an isogeny  
 $\varphi: E \rightarrow E'$



Prover



Verifier



## PROVING KNOWLEDGE OF AN ISOGENY

I know an isogeny  
 $\varphi: E \rightarrow E'$

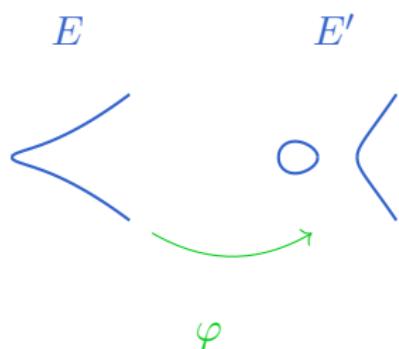


Prover

Really?



Verifier



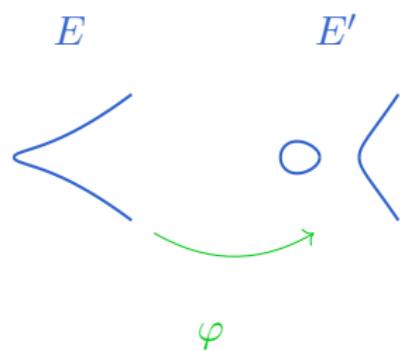
## PROVING KNOWLEDGE OF AN ISOGENY



Prover



Verifier



**NIZK**  
for the relation  $(\varphi, E')$

## PROVING KNOWLEDGE OF AN ISOGENY

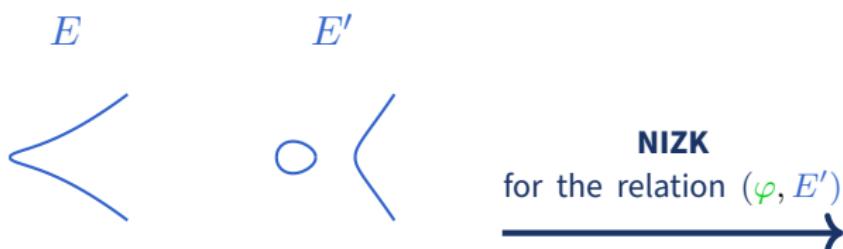


Prover

Oh, I see...



Verifier



**NIZK (Non-interactive zero-knowledge proof):** proving knowledge of the **witness** corresponding to a given **statement**...

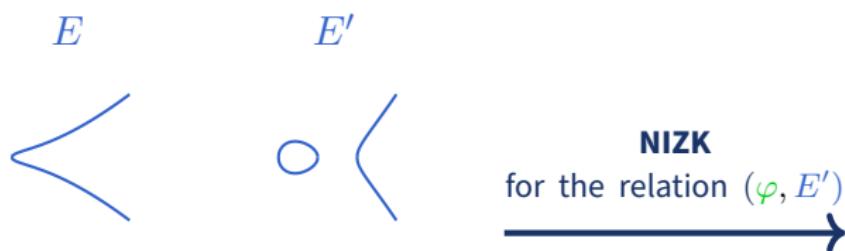
## PROVING KNOWLEDGE OF AN ISOGENY



Prover



Verifier



**NIZK (Non-interactive zero-knowledge proof):** proving knowledge of the **witness** corresponding to a given **statement**...  
...without revealing the **witness**

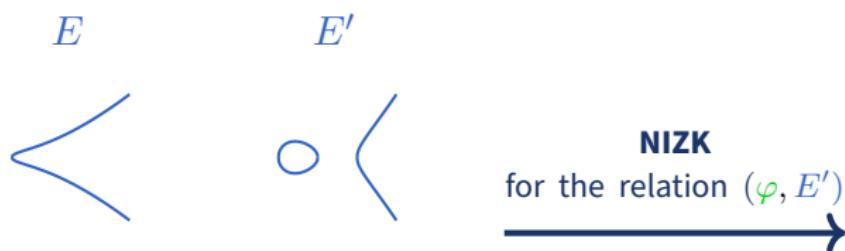
## PROVING KNOWLEDGE OF AN ISOGENY



Prover



Verifier



**NIZK (Non-interactive zero-knowledge proof):** proving knowledge of the **witness** corresponding to a given **statement**...  
...without interacting with the verifier

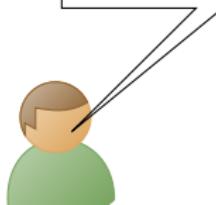
## RANK-1 CONSTRAINT SYSTEMS (R1CS)

A class of relations (**statement**, **witness**) that are ‘easy’ to prove:

I know a vector  $(w_1, \dots, w_n)$  such that

$$\left( \begin{array}{c} \text{A} \end{array} \right) \left( \begin{array}{c} 1 \\ w_1 \\ w_2 \\ \vdots \\ w_n \end{array} \right) \cdot \left( \begin{array}{c} \text{B} \end{array} \right) \left( \begin{array}{c} 1 \\ w_1 \\ w_2 \\ \vdots \\ w_n \end{array} \right) =$$

$$= \left( \begin{array}{c} \text{C} \end{array} \right) \left( \begin{array}{c} 1 \\ w_1 \\ w_2 \\ \vdots \\ w_n \end{array} \right)$$



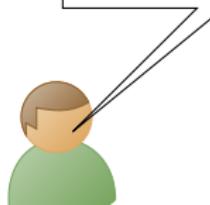
## RANK-1 CONSTRAINT SYSTEMS (R1CS)

A class of relations (**statement**, **witness**) that are ‘easy’ to prove:

I know a vector  $(w_1, \dots, w_n)$  such that

$$\left( \begin{array}{c} \text{A} \end{array} \right) \left( \begin{array}{c} 1 \\ w_1 \\ w_2 \\ \vdots \\ w_n \end{array} \right) \cdot \left( \begin{array}{c} \text{B} \end{array} \right) \left( \begin{array}{c} 1 \\ w_1 \\ w_2 \\ \vdots \\ w_n \end{array} \right) =$$

$$= \left( \begin{array}{c} \text{C} \end{array} \right) \left( \begin{array}{c} 1 \\ w_1 \\ w_2 \\ \vdots \\ w_n \end{array} \right)$$



**Idea (Cong, Lai, Levin):** translate the isogeny relation into an R1CS

## THE ISOGENY RELATION

$E, E' =$  Elliptic curves over  $\mathbb{F}_q$   
 $\varphi: E \rightarrow E'$  = Isogeny of degree  $d$

I know  $\varphi$



## THE ISOGENY RELATION

$E, E' =$  Elliptic curves over  $\mathbb{F}_q$   
 $\varphi: E \rightarrow E'$  = Isogeny of degree  $d$

I know  $\varphi$

=

I know an *efficient representation* of  $\varphi$

## THE ISOGENY RELATION

$E, E' =$  Elliptic curves over  $\mathbb{F}_q$

$\varphi: E \rightarrow E'$  = Isogeny of degree  $d$

I know  $\varphi$

=

I know an *efficient representation* of  $\varphi$

=

I know an algorithm that, on  
input  $P \in E(\mathbb{F}_q)$ , outputs  $\varphi(P)$   
in time  $O(\log q \log d)$

## THE ISOGENY RELATION

$E, E' =$  Elliptic curves over  $\mathbb{F}_q$   
 $\varphi: E \rightarrow E'$  = Isogeny of degree  $d$

I know  $\varphi$

**THEOREM (ROBERT)**

*Every isogeny admits  
an efficient  
representation.*

I know an *efficient representation* of  $\varphi$

I know an algorithm that, on  
input  $P \in E(\mathbb{F}_q)$ , outputs  $\varphi(P)$   
in time  $O(\log q \log d)$

## THE ISOGENY RELATION

$E, E' =$  Elliptic curves over  $\mathbb{F}_q$   
 $\varphi: E \rightarrow E'$  = Isogeny of degree  $d$

I know  $\varphi$

=

I know an *efficient representation* of  $\varphi$

=

I know an algorithm that, on input  $P \in E(\mathbb{F}_q)$ , outputs  $\varphi(P)$  in time  $O(\log q \log d)$

**THEOREM (ROBERT)**

*Every isogeny admits an efficient representation.*

**Our work:** we consider the case  $d = \ell^n$  and  $E, E'$  supersingular.

## WAYS TO REPRESENT AN $\ell^n$ -ISOGENY

- the rational functions defining  $\varphi$
- (Generators of) the kernel of  $\varphi$
- Factor  $\varphi$  as a chain of  $\ell$ -isogenies  $\varphi_n \circ \cdots \circ \varphi_1$
- The  $j$ -invariants of the codomains of  $\varphi_n, \dots, \varphi_1$

## WAYS TO REPRESENT AN $\ell^n$ -ISOGENY

**Example:** over  $\mathbb{F}_{107}$ , the  $3^3$ -isogeny

$$\underbrace{(y^2 = x^3 + x)}_E \xrightarrow{\varphi} \underbrace{(y^2 = x^3 + 55x + 89)}_{E'}$$

- the rational functions defining  $\varphi$
- (Generators of) the kernel of  $\varphi$
- Factor  $\varphi$  as a chain of  $\ell$ -isogenies  $\varphi_n \circ \cdots \circ \varphi_1$
- The  $j$ -invariants of the codomains of  $\varphi_n, \dots, \varphi_1$

## WAYS TO REPRESENT AN $\ell^n$ -ISOGENY

**Example:** over  $\mathbb{F}_{107}$ , the  $3^3$ -isogeny

$$\underbrace{(y^2 = x^3 + x)}_E \xrightarrow{\varphi} \underbrace{(y^2 = x^3 + 55x + 89)}_{E'}$$

- the rational functions defining  $\varphi$

$$\varphi: (x, y) \mapsto \left( \frac{x^{27} + 104x^{26} + \dots}{x^{26} + 104x^{25} + \dots}, \frac{x^{39}y + 49x^{38}y + \dots}{x^{39} + 49x^{38} + \dots} \right)$$

- (Generators of) the kernel of  $\varphi$
- Factor  $\varphi$  as a chain of  $\ell$ -isogenies  $\varphi_n \circ \dots \circ \varphi_1$
- The  $j$ -invariants of the codomains of  $\varphi_n, \dots, \varphi_1$

## WAYS TO REPRESENT AN $\ell^n$ -ISOGENY

**Example:** over  $\mathbb{F}_{107}$ , the  $3^3$ -isogeny

$$\underbrace{(y^2 = x^3 + x)}_E \xrightarrow{\varphi} \underbrace{(y^2 = x^3 + 55x + 89)}_{E'}$$

- the rational functions defining  $\varphi$

$$\varphi: (x, y) \mapsto \left( \frac{x^{27} + 104x^{26} + \dots}{x^{26} + 104x^{25} + \dots}, \frac{x^{39}y + 49x^{38}y + \dots}{x^{39} + 49x^{38} + \dots} \right)$$

- (Generators of) the kernel of  $\varphi$

$$\ker \varphi = \{\text{roots of } x^{26} + 104x^{25} + \dots\} = \langle (3, 43) \rangle$$

- Factor  $\varphi$  as a chain of  $\ell$ -isogenies  $\varphi_n \circ \dots \circ \varphi_1$

- The  $j$ -invariants of the codomains of  $\varphi_n, \dots, \varphi_1$

## WAYS TO REPRESENT AN $\ell^n$ -ISOGENY

**Example:** over  $\mathbb{F}_{107}$ , the  $3^3$ -isogeny

$$\underbrace{(y^2 = x^3 + x)}_E \xrightarrow{\varphi} \underbrace{(y^2 = x^3 + 55x + 89)}_{E'}$$

- the rational functions defining  $\varphi$

$$\varphi: (x, y) \mapsto \left( \frac{x^{27} + 104x^{26} + \dots}{x^{26} + 104x^{25} + \dots}, \frac{x^{39}y + 49x^{38}y + \dots}{x^{39} + 49x^{38} + \dots} \right)$$

- (Generators of) the kernel of  $\varphi$

$$\ker \varphi = \{\text{roots of } x^{26} + 104x^{25} + \dots\} = \langle (3, 43) \rangle$$

- Factor  $\varphi$  as a chain of  $\ell$ -isogenies  $\varphi_n \circ \dots \circ \varphi_1$

$$E \xrightarrow{\varphi_1} (y^2 = x^3 + 89x + 89) \xrightarrow{\varphi_2} (y^2 = x^3 + 86x + 23) \xrightarrow{\varphi_3} E'$$

- The  $j$ -invariants of the codomains of  $\varphi_n, \dots, \varphi_1$

## WAYS TO REPRESENT AN $\ell^n$ -ISOGENY

**Example:** over  $\mathbb{F}_{107}$ , the  $3^3$ -isogeny

$$\underbrace{(y^2 = x^3 + x)}_E \xrightarrow{\varphi} \underbrace{(y^2 = x^3 + 55x + 89)}_{E'}$$

- the rational functions defining  $\varphi$

$$\varphi: (x, y) \mapsto \left( \frac{x^{27} + 104x^{26} + \dots}{x^{26} + 104x^{25} + \dots}, \frac{x^{39}y + 49x^{38}y + \dots}{x^{39} + 49x^{38} + \dots} \right)$$

- (Generators of) the kernel of  $\varphi$

$$\ker \varphi = \{\text{roots of } x^{26} + 104x^{25} + \dots\} = \langle (3, 43) \rangle$$

- Factor  $\varphi$  as a chain of  $\ell$ -isogenies  $\varphi_n \circ \dots \circ \varphi_1$

$$E \xrightarrow{\varphi_1} (y^2 = x^3 + 89x + 89) \xrightarrow{\varphi_2} (y^2 = x^3 + 86x + 23) \xrightarrow{\varphi_3} E'$$

- The  $j$ -invariants of the codomains of  $\varphi_n, \dots, \varphi_1$

$$16 \xrightarrow{\varphi_1} 47 \xrightarrow{\varphi_2} 81 \xrightarrow{\varphi_3} 94$$

## WAYS TO REPRESENT AN $\ell^n$ -ISOGENY

**Example:** over  $\mathbb{F}_{107}$ , the  $3^3$ -isogeny

$$\underbrace{(y^2 = x^3 + x)}_E \xrightarrow{\varphi} \underbrace{(y^2 = x^3 + 55x + 89)}_{E'}$$

- the rational functions defining  $\varphi$

$$\varphi: (x, y) \mapsto \left( \frac{x^{27} + 104x^{26} + \dots}{x^{26} + 104x^{25} + \dots}, \frac{x^{39}y + 49x^{38}y + \dots}{x^{39} + 49x^{38} + \dots} \right)$$

- (Generators of) the kernel of  $\varphi$

$$\ker \varphi = \{\text{roots of } x^{26} + 104x^{25} + \dots\} = \langle (3, 43) \rangle$$

- Factor  $\varphi$  as a chain of  $\ell$ -isogenies  $\varphi_n \circ \dots \circ \varphi_1$

$$E \xrightarrow{\varphi_1} (y^2 = x^3 + 89x + 89) \xrightarrow{\varphi_2} (y^2 = x^3 + 86x + 23) \xrightarrow{\varphi_3} E'$$

- The  $j$ -invariants of the codomains of  $\varphi_n, \dots, \varphi_1$

$$16 \xrightarrow{\varphi_1} 47 \xrightarrow{\varphi_2} 81 \xrightarrow{\varphi_3} 94$$

## FROM ISOGENIES TO R1CS

### Bottom line

The prover can represent his **witness**  $\ell^n$ -isogeny as a chain of  
 $j$ -invariants

$$j_0 = j(E) \xrightarrow{\varphi_1} j_1 \xrightarrow{\varphi_2} j_2 \xrightarrow{\varphi_3} \dots \xrightarrow{\varphi_n} j_n = j(E')$$

## FROM ISOGENIES TO R1CS

### Bottom line

The prover can represent his witness  $\ell^n$ -isogeny as a chain of  $j$ -invariants

$$j_0 = j(E) \xrightarrow{\varphi_1} j_1 \xrightarrow{\varphi_2} j_2 \xrightarrow{\varphi_3} \dots \xrightarrow{\varphi_n} j_n = j(E')$$

**Goal:** plugging  $j_0, j_1, \dots, j_n$  in an R1CS

$$(A) \begin{pmatrix} 1 \\ w_1 \\ \vdots \\ w_n \end{pmatrix} \cdot (B) \begin{pmatrix} 1 \\ w_1 \\ \vdots \\ w_n \end{pmatrix} = (C) \begin{pmatrix} 1 \\ w_1 \\ \vdots \\ w_n \end{pmatrix}$$

## FROM ISOGENIES TO R1CS

### Bottom line

The prover can represent his witness  $\ell^n$ -isogeny as a chain of  $j$ -invariants

$$j_0 = j(E) \xrightarrow{\varphi_1} j_1 \xrightarrow{\varphi_2} j_2 \xrightarrow{\varphi_3} \dots \xrightarrow{\varphi_n} j_n = j(E')$$

**Goal:** plugging  $j_0, j_1, \dots, j_n$  in an R1CS

$$(A) \begin{pmatrix} 1 \\ w_1 \\ \vdots \\ w_n \end{pmatrix} \cdot (B) \begin{pmatrix} 1 \\ w_1 \\ \vdots \\ w_n \end{pmatrix} = (C) \begin{pmatrix} 1 \\ w_1 \\ \vdots \\ w_n \end{pmatrix}$$

To do so, we need equations involving  $j_0, j_1, \dots, j_n$ :

- Cong, Lai, Levin: classical modular polynomials.
- Our work: other modular polynomials.

## CLASSICAL MODULAR POLYNOMIAL FOR $\ell = 2$

$$\begin{aligned}\Phi_2(X, Y) = & X^3 + Y^3 - \underbrace{162000}_{c_2}(X^2 + Y^2) + \underbrace{1488}_{c_4}XY(X + Y) + \\ & - X^2Y^2 + \underbrace{8748000000}_{c_1}(X + Y) + \underbrace{40773375}_{c_3}XY \underbrace{-1574640000000000}_{c_0}\end{aligned}$$

### THEOREM

*There is a 2-isogeny  $j_i \rightarrow j_{i+1}$   $\iff \Phi_2(j_i, j_{i+1}) = 0$*

## CLASSICAL MODULAR POLYNOMIAL FOR $\ell = 2$

$$\begin{aligned}\Phi_2(X, Y) = & X^3 + Y^3 - \underbrace{162000}_{c_2}(X^2 + Y^2) + \underbrace{1488}_{c_4}XY(X + Y) + \\ & - X^2Y^2 + \underbrace{8748000000}_{c_1}(X + Y) + \underbrace{40773375}_{c_3}XY \underbrace{-1574640000000000}_{c_0}\end{aligned}$$

### THEOREM

$$\text{There is a 2-isogeny } j_i \rightarrow j_{i+1} \iff \Phi_2(j_i, j_{i+1}) = 0$$

Therefore, for a  $2^n$ -isogeny...

$$j_0 = j(E) \xrightarrow{\varphi_1} j_1 \xrightarrow{\varphi_2} j_2 \xrightarrow{\varphi_3} \cdots \xrightarrow{\varphi_n} j_n = j(E')$$

## CLASSICAL MODULAR POLYNOMIAL FOR $\ell = 2$

$$\begin{aligned}\Phi_2(X, Y) = & X^3 + Y^3 - \underbrace{162000}_{c_2}(X^2 + Y^2) + \underbrace{1488}_{c_4}XY(X + Y) + \\ & - X^2Y^2 + \underbrace{8748000000}_{c_1}(X + Y) + \underbrace{40773375}_{c_3}XY \underbrace{-1574640000000000}_{c_0}\end{aligned}$$

### THEOREM

$$\text{There is a 2-isogeny } j_i \rightarrow j_{i+1} \iff \Phi_2(j_i, j_{i+1}) = 0$$

Therefore, for a  $2^n$ -isogeny...

$$j_0 = j(E) \xrightarrow{\varphi_1} j_1 \xrightarrow{\varphi_2} j_2 \xrightarrow{\varphi_3} \dots \xrightarrow{\varphi_n} j_n = j(E')$$

$$\Phi_2(j_0, j_1) = 0$$

## CLASSICAL MODULAR POLYNOMIAL FOR $\ell = 2$

$$\begin{aligned}\Phi_2(X, Y) = & X^3 + Y^3 - \underbrace{162000}_{c_2}(X^2 + Y^2) + \underbrace{1488}_{c_4}XY(X + Y) + \\ & - X^2Y^2 + \underbrace{8748000000}_{c_1}(X + Y) + \underbrace{40773375}_{c_3}XY \underbrace{-1574640000000000}_{c_0}\end{aligned}$$

### THEOREM

$$\text{There is a 2-isogeny } j_i \rightarrow j_{i+1} \iff \Phi_2(j_i, j_{i+1}) = 0$$

Therefore, for a  $2^n$ -isogeny...

$$j_0 = j(E) \xrightarrow{\varphi_1} j_1 \xrightarrow{\varphi_2} j_2 \xrightarrow{\varphi_3} \dots \xrightarrow{\varphi_n} j_n = j(E')$$

$$\Phi_2(j_1, j_2) = 0$$

## CLASSICAL MODULAR POLYNOMIAL FOR $\ell = 2$

$$\begin{aligned}\Phi_2(X, Y) = & X^3 + Y^3 - \underbrace{162000}_{c_2}(X^2 + Y^2) + \underbrace{1488}_{c_4}XY(X + Y) + \\ & - X^2Y^2 + \underbrace{8748000000}_{c_1}(X + Y) + \underbrace{40773375}_{c_3}XY \underbrace{-1574640000000000}_{c_0}\end{aligned}$$

### THEOREM

$$\text{There is a 2-isogeny } j_i \rightarrow j_{i+1} \iff \Phi_2(j_i, j_{i+1}) = 0$$

Therefore, for a  $2^n$ -isogeny...

$$j_0 = j(E) \xrightarrow{\varphi_1} j_1 \xrightarrow{\varphi_2} j_2 \xrightarrow{\varphi_3} \dots \xrightarrow{\varphi_n} j_n = j(E')$$

$$\Phi_2(j_2, \dots) = 0$$

## CLASSICAL MODULAR POLYNOMIAL FOR $\ell = 2$

$$\begin{aligned}\Phi_2(X, Y) = & X^3 + Y^3 - \underbrace{162000}_{c_2}(X^2 + Y^2) + \underbrace{1488}_{c_4}XY(X + Y) + \\ & - X^2Y^2 + \underbrace{8748000000}_{c_1}(X + Y) + \underbrace{40773375}_{c_3}XY \underbrace{-1574640000000000}_{c_0}\end{aligned}$$

### THEOREM

$$\text{There is a 2-isogeny } j_i \rightarrow j_{i+1} \iff \Phi_2(j_i, j_{i+1}) = 0$$

Therefore, for a  $2^n$ -isogeny...

$$j_0 = j(E) \xrightarrow{\varphi_1} j_1 \xrightarrow{\varphi_2} j_2 \xrightarrow{\varphi_3} \cdots \xrightarrow{\varphi_n} j_n = j(E')$$

$$\Phi_2(j_{n-1}, j_n) = 0$$

## CLASSICAL MODULAR POLYNOMIAL FOR ANY $\ell$

$\Phi_\ell(X, Y) = X^{\ell+1} + Y^{\ell+1} + \text{polynomial symmetric in } X \text{ and } Y$   
with very large integer coefficients

### THEOREM

*There is an  $\ell$ -isogeny  $j_i \rightarrow j_{i+1}$*        $\iff$        $\Phi_\ell(j_i, j_{i+1}) = 0$

Therefore, for an  $\ell^n$ -isogeny...

$$j_0 = j(E) \xrightarrow{\varphi_1} j_1 \xrightarrow{\varphi_2} j_2 \xrightarrow{\varphi_3} \cdots \xrightarrow{\varphi_n} j_n = j(E')$$

$$\Phi_\ell(j_{n-1}, j_n) = 0$$

$$\begin{aligned} (\mathbf{A}) \begin{pmatrix} 1 \\ w_1 \\ \vdots \\ w_n \end{pmatrix} \cdot (\mathbf{B}) \begin{pmatrix} 1 \\ w_1 \\ \vdots \\ w_n \end{pmatrix} &= \\ &= (\mathbf{C}) \begin{pmatrix} 1 \\ w_1 \\ \vdots \\ w_n \end{pmatrix} \end{aligned}$$

$$\begin{pmatrix} 0 & 1 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 1 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & -1 \end{pmatrix} \begin{pmatrix} 1 \\ j_i \\ j_{i+1} \\ j_i^2 \\ j_{i+1}^2 \\ j_i^3 \\ j_{i+1}^3 \\ j_i j_{i+1} \end{pmatrix} \cdot \begin{pmatrix} 0 & 1 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 & 0 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 & 0 & 0 & 0 & 0 \\ 0 & c_4 & c_4 & 0 & 0 & 0 & 0 & -1 \end{pmatrix} \begin{pmatrix} 1 \\ j_i \\ j_{i+1} \\ j_i^2 \\ j_{i+1}^2 \\ j_i^3 \\ j_{i+1}^3 \\ j_i j_{i+1} \end{pmatrix} = \\
 = \begin{pmatrix} 0 & 0 & 0 & 1 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 1 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 \\ c_0 & c_1 & c_1 & c_2 & c_2 & 1 & 1 & c_3 \end{pmatrix} \begin{pmatrix} 1 \\ j_i \\ j_{i+1} \\ j_i^2 \\ j_{i+1}^2 \\ j_i^3 \\ j_{i+1}^3 \\ j_i j_{i+1} \end{pmatrix}$$

...for each  $i$  in  $\{0, \dots, n - 1\}$

## OUR WORK: FROM CLASSICAL TO CANONICAL MODULAR POLYNOMIALS

For  $\ell \in \{2, 3, 5, 7, 13\}$ , we consider the *canonical polynomials* constructed by Müller

$$\Phi_2^c(f, j) = f^3 + 48f^2 + 768f + 4096 - f \cdot j$$

$$\Phi_3^c(f, j) = f^4 + 36f^3 + 270f^2 + 756f + 729 - f \cdot j$$

$$\begin{aligned} \Phi_5^c(f, j) = & f^6 + 30f^5 + 315f^4 + 1300f^3 + 1575f^2 + 750f + \\ & + 125 - f \cdot j \end{aligned}$$

$$\begin{aligned} \Phi_7^c(f, j) = & f^8 + 28f^7 + 322f^6 + 1904f^5 + 5915f^4 + 8624f^3 + \\ & + 4018f^2 + 748f + 49 - f \cdot j \end{aligned}$$

$$\begin{aligned} \Phi_{13}^c(f, j) = & f^{14} + 26f^{13} + 325f^{12} + 2548f^{11} + 13832f^{10} + \\ & + 54340f^9 + 157118f^8 + 333580f^7 + 509366f^6 + \\ & + 534820f^5 + 354536f^4 + 124852f^3 + 15145f^2 + \\ & + 746f + 13 - f \cdot j \end{aligned}$$

## OUR WORK: FROM CLASSICAL TO CANONICAL MODULAR POLYNOMIALS

For  $\ell \in \{2, 3, 5, 7, 13\}$ , we consider the *canonical polynomials* constructed by Müller

$$\Phi_3^c(f, j) = f^4 + 36f^3 + 270f^2 + 756f + 729 - f \cdot j$$

For comparison, the 3rd classical modular polynomial is

$$\begin{aligned}\Phi_3(X, Y) = & -X^3Y^3 + X^4 + Y^4 + 2232(X^3Y^2 + X^2Y^3) \\ & - 1069956(X^3Y + XY^3) + 36864000(X^3 + Y^3) \\ & + 2587918086X^2Y^2 + 8900222976000(X^2Y + XY^2) \\ & + 452984832000000(X^2 + Y^2) - 770845966336000000XY \\ & + 1855425871872000000000(X + Y)\end{aligned}$$

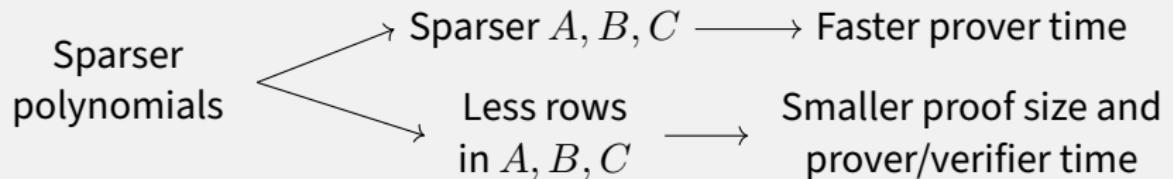
## OUR WORK: FROM CLASSICAL TO CANONICAL MODULAR POLYNOMIALS

For  $\ell \in \{2, 3, 5, 7, 13\}$ , we consider the *canonical polynomials* constructed by Müller

$$\Phi_3^c(f, j) = f^4 + 36f^3 + 270f^2 + 756f + 729 - f \cdot j$$

For comparison, the 3rd classical modular polynomial is

$$\begin{aligned}\Phi_3(X, Y) = & -X^3Y^3 + X^4 + Y^4 + 2232(X^3Y^2 + X^2Y^3) \\ & - 1069956(X^3Y + XY^3) + 36864000(X^3 + Y^3) \\ & + 2587918086X^2Y^2 + 8900222976000(X^2Y + XY^2) \\ & + 452984832000000(X^2 + Y^2) - 770845966336000000XY \\ & + 1855425871872000000000(X + Y)\end{aligned}$$



## OUR MAIN RESULTS ON CANONICAL MODULAR POLYNOMIALS

- How does  $\Phi_\ell^c$  encode isogenies?

### THEOREM

Let  $\ell \in \{2, 3, 5, 7, 13\}$  and  $s = \frac{12}{\gcd(12, \ell-1)}$ .

*There is an  $\ell$ -isogeny  $j_i \rightarrow j_{i+1}$*

$$\Updownarrow$$

*There exists  $f_i \in \overline{\mathbb{F}_p}^\times$  s.t.  $\begin{cases} \Phi_\ell^c(f_i, j_i) = 0 \\ \Phi_\ell^c(\ell^s / f_i, j_{i+1}) = 0 \end{cases}$*

## OUR MAIN RESULTS ON CANONICAL MODULAR POLYNOMIALS

- How does  $\Phi_\ell^c$  encode isogenies?

### THEOREM

Let  $\ell \in \{2, 3, 5, 7, 13\}$  and  $s = \frac{12}{\gcd(12, \ell-1)}$ .

*There is an  $\ell$ -isogeny  $j_i \rightarrow j_{i+1}$*

$$\Updownarrow$$

*There exists  $f_i \in \overline{\mathbb{F}_p}^\times$  s.t.  $\begin{cases} \Phi_\ell^c(f_i, j_i) = 0 \\ \Phi_\ell^c(\ell^s / f_i, j_{i+1}) = 0 \end{cases}$*

$$j_0 = j(E) \xrightarrow{\varphi_1} j_1 \xrightarrow{\varphi_2} j_2 \xrightarrow{\varphi_3} \cdots \xrightarrow{\varphi_n} j_n = j(E')$$

$$\begin{cases} \Phi_\ell^c(f_0, j_0) = 0 \\ \Phi_\ell^c(\ell^s / f_0, j_1) = 0 \end{cases}$$

## OUR MAIN RESULTS ON CANONICAL MODULAR POLYNOMIALS

- How does  $\Phi_\ell^c$  encode isogenies?

### THEOREM

Let  $\ell \in \{2, 3, 5, 7, 13\}$  and  $s = \frac{12}{\gcd(12, \ell-1)}$ .

*There is an  $\ell$ -isogeny  $j_i \rightarrow j_{i+1}$*

$$\Updownarrow$$

*There exists  $f_i \in \overline{\mathbb{F}_p}^\times$  s.t.  $\begin{cases} \Phi_\ell^c(f_i, j_i) = 0 \\ \Phi_\ell^c(\ell^s / f_i, j_{i+1}) = 0 \end{cases}$*

$$j_0 = j(E) \xrightarrow{\varphi_1} j_1 \xrightarrow{\varphi_2} j_2 \xrightarrow{\varphi_3} \cdots \xrightarrow{\varphi_n} j_n = j(E')$$

$$\begin{cases} \Phi_\ell^c(f_1, j_1) = 0 \\ \Phi_\ell^c(\ell^s / f_1, j_2) = 0 \end{cases}$$

## OUR MAIN RESULTS ON CANONICAL MODULAR POLYNOMIALS

- How does  $\Phi_\ell^c$  encode isogenies?

### THEOREM

Let  $\ell \in \{2, 3, 5, 7, 13\}$  and  $s = \frac{12}{\gcd(12, \ell-1)}$ .

*There is an  $\ell$ -isogeny  $j_i \rightarrow j_{i+1}$*

$$\Updownarrow$$

*There exists  $f_i \in \overline{\mathbb{F}_p}^\times$  s.t.  $\begin{cases} \Phi_\ell^c(f_i, j_i) = 0 \\ \Phi_\ell^c(\ell^s / f_i, j_{i+1}) = 0 \end{cases}$*

$$j_0 = j(E) \xrightarrow{\varphi_1} j_1 \xrightarrow{\varphi_2} j_2 \xrightarrow{\varphi_3} \cdots \xrightarrow{\varphi_n} j_n = j(E')$$

$$\begin{cases} \Phi_\ell^c(f_2, j_2) = 0 \\ \Phi_\ell^c(\ell^s / f_2, \dots) = 0 \end{cases}$$

## OUR MAIN RESULTS ON CANONICAL MODULAR POLYNOMIALS

- How does  $\Phi_\ell^c$  encode isogenies?

### THEOREM

Let  $\ell \in \{2, 3, 5, 7, 13\}$  and  $s = \frac{12}{\gcd(12, \ell-1)}$ .

*There is an  $\ell$ -isogeny  $j_i \rightarrow j_{i+1}$*

$$\Updownarrow$$

*There exists  $f_i \in \overline{\mathbb{F}_p}^\times$  s.t.  $\begin{cases} \Phi_\ell^c(f_i, j_i) = 0 \\ \Phi_\ell^c(\ell^s / f_i, j_{i+1}) = 0 \end{cases}$*

$$j_0 = j(E) \xrightarrow{\varphi_1} j_1 \xrightarrow{\varphi_2} j_2 \xrightarrow{\varphi_3} \cdots \xrightarrow{\varphi_n} j_n = j(E')$$

$$\begin{cases} \Phi_\ell^c(f_{n-1}, j_{n-1}) = 0 \\ \Phi_\ell^c(\ell^s / f_{n-1}, j_n) = 0 \end{cases}$$

## OUR MAIN RESULTS ON CANONICAL MODULAR POLYNOMIALS

- How does  $\Phi_\ell^c$  encode isogenies?

### THEOREM

Let  $\ell \in \{2, 3, 5, 7, 13\}$  and  $s = \frac{12}{\gcd(12, \ell-1)}$ .

*There is an  $\ell$ -isogeny  $j_i \rightarrow j_{i+1}$*

$$\Updownarrow$$

*There exists  $f_i \in \overline{\mathbb{F}_p}^\times$  s.t.  $\begin{cases} \Phi_\ell^c(f_i, j_i) = 0 \\ \Phi_\ell^c(\ell^s / f_i, j_{i+1}) = 0 \end{cases}$*

- Where does  $f_i$  live?

### THEOREM

Let  $\ell \in \{2, 3, 5, 7, 13\}$ . If  $j_i$  is supersingular, then  $\Phi_\ell^c(f, j_i)$  splits over  $\mathbb{F}_{p^2}$ .

## R1CS: CLASSICAL VS CANONICAL

Cong, Lai, Levin:  $\ell = 2$  with classical modular polynomials

$$\begin{pmatrix} 0 & 1 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 1 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & -1 \end{pmatrix} \begin{pmatrix} 1 \\ j_i \\ j_{i+1} \\ j_i^2 \\ j_{i+1}^2 \\ j_i^3 \\ j_{i+1}^3 \\ j_i j_{i+1} \end{pmatrix} \cdot \begin{pmatrix} 0 & 1 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 & 0 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 & 0 & 0 & 0 & 0 \\ 0 & c_4 & c_4 & 0 & 0 & 0 & 0 & -1 \end{pmatrix} \begin{pmatrix} 1 \\ j_i \\ j_{i+1} \\ j_i^2 \\ j_{i+1}^2 \\ j_i^3 \\ j_{i+1}^3 \\ j_i j_{i+1} \end{pmatrix} = \\
 = \begin{pmatrix} 0 & 0 & 0 & 1 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 1 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 \\ c_0 & c_1 & c_1 & c_2 & c_2 & 1 & 1 & c_3 \end{pmatrix} \begin{pmatrix} 1 \\ j_i \\ j_{i+1} \\ j_i^2 \\ j_{i+1}^2 \\ j_i^3 \\ j_{i+1}^3 \\ j_i j_{i+1} \end{pmatrix}$$

...for each  $i$  in  $\{0, \dots, n - 1\}$

## R1CS: CLASSICAL VS CANONICAL

**Our work:**  $\ell = 2$  with canonical modular polynomials

$$\begin{pmatrix} 0 & 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 & 0 \end{pmatrix} \begin{pmatrix} 1 \\ f_i \\ f_i^2 \\ j_i - c'_1 \\ j_{i+1} - c'_1 \end{pmatrix} \cdot \begin{pmatrix} 0 & 1 & 0 & 0 & 0 \\ 0 & c'_2 & c'_3 & -1 & 0 \\ 0 & c''_3 & 0 & 0 & -1 \end{pmatrix} \begin{pmatrix} 1 \\ f_i \\ f_i^2 \\ j_i - c'_1 \\ j_{i+1} - c'_1 \end{pmatrix} = \\ = \begin{pmatrix} 0 & 0 & 1 & 0 & 0 \\ -c'_0 & 0 & 0 & 0 & 0 \\ -c''_0 & -c''_1 & 0 & 0 & 0 \end{pmatrix} \begin{pmatrix} 1 \\ f_i \\ f_i^2 \\ j_i - c'_1 \\ j_{i+1} - c'_1 \end{pmatrix}$$

...for each  $i$  in  $\{0, \dots, n-1\}$

(and  $c'_k, c''_k$  obtained from the coefficients of  $\Phi_2^c(f, X)$ )

## OUR WORK (IN PROGRESS): ATKIN MODULAR POLYNOMIALS

For  $\ell \in \{2, 3, 5, 7, 11, 13, 17, 19, 23, 29, 31, 41, 47, 59, 71\}$ , we consider the Atkin modular polynomials:

$$\begin{aligned}\Phi_2^A(f, j) = & j^2 - j(f^2 + f - 7256) + f^3 + 744f^2 + 184512f + \\& + 15252992\end{aligned}$$

$$\begin{aligned}\Phi_3^A(f, j) = & j^2 - j(f^3 - 2348f - 24528) + f^4 + 744f^3 + \\& + 193752f^2 + 19712160f + 538141968\end{aligned}$$

⋮

$$\Phi_\ell^A(f, j) = j^2 - j \cdot \left( f^\ell + \sum_{i=0}^{\ell-1} a_{\ell,i} f^i \right) + f^{\ell+1} + \sum_{i=0}^{\ell} b_{\ell,i} f^i$$

## OUR WORK (IN PROGRESS): ATKIN MODULAR POLYNOMIALS

For  $\ell \in \{2, 3, 5, 7, 11, 13, 17, 19, 23, 29, 31, 41, 47, 59, 71\}$ , we consider the Atkin modular polynomials:

$$\begin{aligned}\Phi_2^A(f, j) = & j^2 - j(f^2 + f - 7256) + f^3 + 744f^2 + 184512f + \\ & + 15252992\end{aligned}$$

$$\begin{aligned}\Phi_3^A(f, j) = & j^2 - j(f^3 - 2348f - 24528) + f^4 + 744f^3 + \\ & + 193752f^2 + 19712160f + 538141968\end{aligned}$$

⋮

$$\Phi_\ell^A(f, j) = j^2 - j \cdot \left( f^\ell + \sum_{i=0}^{\ell-1} a_{\ell,i} f^i \right) + f^{\ell+1} + \sum_{i=0}^{\ell} b_{\ell,i} f^i$$

We also consider the difference quotient polynomial

$$\delta_\ell(f, j_0, j_1) = \frac{\Phi_\ell^A(f, j_1) - \Phi_\ell^A(f, j_0)}{j_1 - j_0} = j_0 + j_1 - \left( f^\ell + \sum_{i=0}^{\ell-1} a_{\ell,i} f^i \right)$$

## OUR RESULTS ON ATKIN MODULAR POLYNOMIALS

- How does  $\Phi_\ell^A$  encode isogenies?

### THEOREM

Let  $\ell \in \{2, 3, 5, 7, 11, 13, 17, 19, 23, 29, 31, 41, 47, 59, 71\}$ .

*There is an  $\ell$ -isogeny  $j_i \rightarrow j_{i+1}$*

$$\Updownarrow$$

*There exists  $f_i \in \overline{\mathbb{F}_p}^\times$  s.t.  $\begin{cases} \Phi_\ell^A(f_i, j_i) = 0 \\ \delta_\ell(f_i, j_i, j_{i+1}) = 0 \end{cases}$*

## OUR RESULTS ON ATKIN MODULAR POLYNOMIALS

- How does  $\Phi_\ell^A$  encode isogenies?

### THEOREM

Let  $\ell \in \{2, 3, 5, 7, 11, 13, 17, 19, 23, 29, 31, 41, 47, 59, 71\}$ .

*There is an  $\ell$ -isogeny  $j_i \rightarrow j_{i+1}$*

$\Updownarrow$

*There exists  $f_i \in \overline{\mathbb{F}_p}^\times$  s.t.  $\begin{cases} \Phi_\ell^A(f_i, j_i) = 0 \\ \delta_\ell(f_i, j_i, j_{i+1}) = 0 \end{cases}$*

$$j_0 = j(E) \xrightarrow{\varphi_1} j_1 \xrightarrow{\varphi_2} j_2 \xrightarrow{\varphi_3} \dots \xrightarrow{\varphi_n} j_n = j(E')$$

$$\begin{cases} \Phi_\ell^A(f_0, j_0) = 0 \\ \delta_\ell(f_0, j_0, j_1) = 0 \end{cases}$$

## OUR RESULTS ON ATKIN MODULAR POLYNOMIALS

- How does  $\Phi_\ell^A$  encode isogenies?

### THEOREM

Let  $\ell \in \{2, 3, 5, 7, 11, 13, 17, 19, 23, 29, 31, 41, 47, 59, 71\}$ .

*There is an  $\ell$ -isogeny  $j_i \rightarrow j_{i+1}$*

$$\Updownarrow$$

*There exists  $f_i \in \overline{\mathbb{F}_p}^\times$  s.t.  $\begin{cases} \Phi_\ell^A(f_i, j_i) = 0 \\ \delta_\ell(f_i, j_i, j_{i+1}) = 0 \end{cases}$*

$$j_0 = j(E) \xrightarrow{\varphi_1} j_1 \xrightarrow{\varphi_2} j_2 \xrightarrow{\varphi_3} \dots \xrightarrow{\varphi_n} j_n = j(E')$$

$$\begin{cases} \Phi_\ell^A(f_1, j_1) = 0 \\ \delta_\ell(f_1, j_1, j_2) = 0 \end{cases}$$

## OUR RESULTS ON ATKIN MODULAR POLYNOMIALS

- How does  $\Phi_\ell^A$  encode isogenies?

### THEOREM

Let  $\ell \in \{2, 3, 5, 7, 11, 13, 17, 19, 23, 29, 31, 41, 47, 59, 71\}$ .

*There is an  $\ell$ -isogeny  $j_i \rightarrow j_{i+1}$*

$$\Updownarrow$$

*There exists  $f_i \in \overline{\mathbb{F}_p}^\times$  s.t.  $\begin{cases} \Phi_\ell^A(f_i, j_i) = 0 \\ \delta_\ell(f_i, j_i, j_{i+1}) = 0 \end{cases}$*

$$j_0 = j(E) \xrightarrow{\varphi_1} j_1 \xrightarrow{\varphi_2} j_2 \xrightarrow{\varphi_3} \dots \xrightarrow{\varphi_n} j_n = j(E')$$

$$\begin{cases} \Phi_\ell^A(f_2, j_2) = 0 \\ \delta_\ell(f_2, j_2, \dots) = 0 \end{cases}$$

## OUR RESULTS ON ATKIN MODULAR POLYNOMIALS

- How does  $\Phi_\ell^A$  encode isogenies?

### THEOREM

Let  $\ell \in \{2, 3, 5, 7, 11, 13, 17, 19, 23, 29, 31, 41, 47, 59, 71\}$ .

*There is an  $\ell$ -isogeny  $j_i \rightarrow j_{i+1}$*

$\Updownarrow$

*There exists  $f_i \in \overline{\mathbb{F}_p}^\times$  s.t.* 
$$\begin{cases} \Phi_\ell^A(f_i, j_i) = 0 \\ \delta_\ell(f_i, j_i, j_{i+1}) = 0 \end{cases}$$

$$j_0 = j(E) \xrightarrow{\varphi_1} j_1 \xrightarrow{\varphi_2} j_2 \xrightarrow{\varphi_3} \cdots \xrightarrow{\varphi_n} j_n = j(E')$$

$$\begin{cases} \Phi_\ell^A(f_{n-1}, j_{n-1}) = 0 \\ \delta_\ell(f_{n-1}, j_{n-1}, j_n) = 0 \end{cases}$$

## OUR RESULTS ON ATKIN MODULAR POLYNOMIALS

- How does  $\Phi_\ell^A$  encode isogenies?

### THEOREM

Let  $\ell \in \{2, 3, 5, 7, 11, 13, 17, 19, 23, 29, 31, 41, 47, 59, 71\}$ .

*There is an  $\ell$ -isogeny  $j_i \rightarrow j_{i+1}$*

$$\Updownarrow$$

*There exists  $f_i \in \overline{\mathbb{F}_p}^\times$  s.t.  $\begin{cases} \Phi_\ell^A(f_i, j_i) = 0 \\ \delta_\ell(f_i, j_i, j_{i+1}) = 0 \end{cases}$*

- Where does  $f_i$  live?

**Conjecture:** If  $j_i$  is supersingular, then  $\Phi_\ell^A(f, j_i)$  splits over  $\mathbb{F}_{p^2}$

## MODULAR CURVES: THE ELEPHANT IN THE ROOM

$$\{\text{Elliptic curves}\}/\simeq \longleftrightarrow X(1)$$

$$\{\ell\text{-isogenies}\}/\sim \longleftrightarrow X_0(\ell)$$

## MODULAR CURVES: THE ELEPHANT IN THE ROOM

$\{\text{Elliptic curves}\}/\simeq \longleftrightarrow X(1)$



$\{\ell\text{-isogenies}\}/\sim \longleftrightarrow X_0(\ell)$

## MODULAR CURVES: THE ELEPHANT IN THE ROOM

$$\{\text{Elliptic curves}\}/\sim \longleftrightarrow X(1)$$

$$\begin{array}{c} \left\{ \{ \varphi, \hat{\varphi} \} \mid \right. \\ \left. \varphi \text{ is an } \ell\text{-isogeny} \right\} / \sim \longleftrightarrow X_0^+(\ell) \\ \uparrow \qquad \qquad \qquad \uparrow \\ \{\ell\text{-isogenies}\} / \sim \longleftrightarrow X_0(\ell) \end{array}$$

## MODULAR CURVES: THE ELEPHANT IN THE ROOM

$$\{\text{Elliptic curves}\}/\sim \longleftrightarrow X(1) \quad \text{Classical mod. poly}$$

$$\begin{array}{c} \left\{ \{ \varphi, \hat{\varphi} \} \mid \right. \\ \left. \varphi \text{ is an } \ell\text{-isogeny} \right\} / \sim \longleftrightarrow X_0^+(\ell) \\ \uparrow \qquad \qquad \qquad \uparrow \\ \{\ell\text{-isogenies}\} / \sim \longleftrightarrow X_0(\ell) \end{array}$$

## MODULAR CURVES: THE ELEPHANT IN THE ROOM

$$\{\text{Elliptic curves}\}/\sim \longleftrightarrow X(1) \quad \text{Classical mod. poly}$$

$$\begin{array}{c} \left\{ \{ \varphi, \hat{\varphi} \} \mid \right. \\ \left. \varphi \text{ is an } \ell\text{-isogeny} \right\} / \sim \longleftrightarrow X_0^+(\ell) \\ \uparrow \qquad \qquad \qquad \uparrow \\ \{\ell\text{-isogenies}\}/\sim \longleftrightarrow X_0(\ell) \quad \text{Canonical mod. poly} \end{array}$$

## MODULAR CURVES: THE ELEPHANT IN THE ROOM

$$\{\text{Elliptic curves}\}/\sim \longleftrightarrow X(1) \quad \text{Classical mod. poly}$$

$$\begin{array}{c} \left\{ \{\varphi, \hat{\varphi}\} \mid \right. \\ \left. \varphi \text{ is an } \ell\text{-isogeny} \right\} / \sim \longleftrightarrow X_0^+(\ell) \quad \text{Atkin mod. poly} \\ \uparrow \qquad \qquad \qquad \uparrow \\ \{\ell\text{-isogenies}\} / \sim \longleftrightarrow X_0(\ell) \quad \text{Canonical mod. poly} \end{array}$$

## MODULAR CURVES: THE ELEPHANT IN THE ROOM

$$\{\text{Elliptic curves}\}/\sim \longleftrightarrow X(1) \quad \text{Classical mod. poly}$$

$$\begin{array}{c} \left\{ \left\{ \varphi, \hat{\varphi} \right\} \mid \right. \\ \left. \varphi \text{ is an } \ell\text{-isogeny} \right\} / \sim \longleftrightarrow X_0^+(\ell) \quad \text{Atkin mod. poly} \\ \uparrow \qquad \qquad \qquad \uparrow \\ \{\ell\text{-isogenies}\}/\sim \longleftrightarrow X_0(\ell) \quad \text{Canonical mod. poly} \end{array}$$

**Our results apply exactly when  
 $X_0(\ell)$  (resp.  $X_0^+(\ell)$ ) has genus 0.**



**THANK YOU FOR YOUR ATTENTION!**

## ESSENTIAL BIBLIOGRAPHY

- [CLL23] K. Cong, Y.-F. Lai, and S. Levin. “Efficient Isogeny Proofs Using Generic Techniques”. In: *Applied Cryptography and Network Security*. Cham: Springer Nature Switzerland, 2023, pp. 248–275.
- [Hol+] T. den Hollander et al. *More Efficient Isogeny Proofs of Knowledge via Canonical Modular Polynomials*. To appear at CRYPTO 2025. URL: <https://eprint.iacr.org/2024/1738>.
- [Mü95] V. Müller. “Ein Algorithmus zur Bestimmung der Punktanzahl elliptischer Kurven über endlichen Körpern der Charakteristik größer drei”. PhD thesis. Universität des Saarlandes, 1995.
- [Rob] D. Robert. *On the efficient representation of isogenies (a survey)*. URL: <https://eprint.iacr.org/2024/1071>.

## DATA COMPARISON

$\ell$	Rows		Variables		Non-zero entries	
	CLL	Ours	CLL	Ours	CLL	Ours
2	$4\lambda + 2$	$3\lambda$	$4\lambda + 3$	$3\lambda + 1$	$21\lambda + 6$	$13\lambda$
3		$2.524\lambda$		$2.524\lambda + 1$		$11.357\lambda$
5		$2.584\lambda$		$2.584\lambda + 1$		$12.059\lambda$
7		$2.493\lambda$		$2.493\lambda + 1$		$12.467\lambda$
13		$2.702\lambda$		$2.702\lambda + 1$		$15.133\lambda$

Table: Parameters of the R1CS for an  $\ell$ -smooth isogeny of degree  $\geq 2^\lambda$

## WAYS TO REPRESENT AN $\ell^n$ -ISOGENY – EXAMPLE OVER $\mathbb{F}_{41^2}$

- the rational functions defining  $\varphi$
- (Generators of) the kernel of  $\varphi$  + Vélu's formulae
- Factor  $\varphi$  as a chain of  $\ell$ -isogenies  $\varphi_n \circ \cdots \circ \varphi_1$
- The  $j$ -invariants of the codomains of  $\varphi_n, \dots, \varphi_1$

## WAYS TO REPRESENT AN $\ell^n$ -ISOGENY – EXAMPLE OVER $\mathbb{F}_{41^2}$

**Example:** over  $\mathbb{F}_{41^2}$ , the  $5^3$ -isogeny

$$\underbrace{(y^2 = x^3 + 1)}_E \xrightarrow{\varphi} \underbrace{(y^2 = x^3 + (16a + 10)x + (3a + 34))}_{E'}$$

with  $(a - 22)^2 \equiv 27 \pmod{41}$

- the rational functions defining  $\varphi$
- (Generators of) the kernel of  $\varphi$  + Vélu's formulae
- Factor  $\varphi$  as a chain of  $\ell$ -isogenies  $\varphi_n \circ \cdots \circ \varphi_1$
- The  $j$ -invariants of the codomains of  $\varphi_n, \dots, \varphi_1$

## WAYS TO REPRESENT AN $\ell^n$ -ISOGENY – EXAMPLE OVER $\mathbb{F}_{41^2}$

**Example:** over  $\mathbb{F}_{41^2}$ , the  $5^3$ -isogeny

$$\underbrace{(y^2 = x^3 + 1)}_E \xrightarrow{\varphi} \underbrace{(y^2 = x^3 + (16a + 10)x + (3a + 34))}_{E'} \quad \text{with } (a - 22)^2 \equiv 27 \pmod{41}$$

- the rational functions defining  $\varphi$

$$\varphi: (x, y) \mapsto \left( \frac{4x^{125} + (7a + 7)x^{124} + \dots}{10x^{124} + (38a + 38)x^{123} + \dots}, \frac{21x^{186}y + \dots}{31x^{186} + \dots} \right)$$

- (Generators of) the kernel of  $\varphi$  + Vélu's formulae

- Factor  $\varphi$  as a chain of  $\ell$ -isogenies  $\varphi_n \circ \dots \circ \varphi_1$

- The  $j$ -invariants of the codomains of  $\varphi_n, \dots, \varphi_1$

## WAYS TO REPRESENT AN $\ell^n$ -ISOGENY – EXAMPLE OVER $\mathbb{F}_{41^2}$

**Example:** over  $\mathbb{F}_{41^2}$ , the  $5^3$ -isogeny

$$\underbrace{(y^2 = x^3 + 1)}_E \xrightarrow{\varphi} \underbrace{(y^2 = x^3 + (16a + 10)x + (3a + 34))}_{E'} \quad \text{with } (a - 22)^2 \equiv 27 \pmod{41}$$

- the rational functions defining  $\varphi$

$$\varphi: (x, y) \mapsto \left( \frac{4x^{125} + (7a + 7)x^{124} + \dots}{10x^{124} + (38a + 38)x^{123} + \dots}, \frac{21x^{186}y + \dots}{31x^{186} + \dots} \right)$$

- (Generators of) the kernel of  $\varphi$  + Vélu's formulae

$$\ker \varphi = \{ \text{roots of } x^{62} + (6a + 6)x^{61} + \dots \} \subseteq \mathbb{F}_{41^{50}}$$

- Factor  $\varphi$  as a chain of  $\ell$ -isogenies  $\varphi_n \circ \dots \circ \varphi_1$

- The  $j$ -invariants of the codomains of  $\varphi_n, \dots, \varphi_1$

## WAYS TO REPRESENT AN $\ell^n$ -ISOGENY – EXAMPLE OVER $\mathbb{F}_{41^2}$

**Example:** over  $\mathbb{F}_{41^2}$ , the  $5^3$ -isogeny

$$\underbrace{(y^2 = x^3 + 1)}_E \xrightarrow{\varphi} \underbrace{(y^2 = x^3 + (16a + 10)x + (3a + 34))}_{E'} \quad \text{with } (a - 22)^2 \equiv 27 \pmod{41}$$

- the rational functions defining  $\varphi$

$$\varphi: (x, y) \mapsto \left( \frac{4x^{125} + (7a + 7)x^{124} + \dots}{10x^{124} + (38a + 38)x^{123} + \dots}, \frac{21x^{186}y + \dots}{31x^{186} + \dots} \right)$$

- (Generators of) the kernel of  $\varphi$  + Vélu's formulae

$$\ker \varphi = \{\text{roots of } x^{62} + (6a + 6)x^{61} + \dots\} \subseteq \mathbb{F}_{41^{50}}$$

- Factor  $\varphi$  as a chain of  $\ell$ -isogenies  $\varphi_n \circ \dots \circ \varphi_1$

$$\begin{aligned} E &\xrightarrow{\varphi_1} (y^2 = x^3 + (21a + 23)x + 1) \\ &\xrightarrow{\varphi_2} (y^2 = x^3 + (34a + 11)x + (21a + 33)) \xrightarrow{\varphi_3} E' \end{aligned}$$

- The  $j$ -invariants of the codomains of  $\varphi_n, \dots, \varphi_1$

## WAYS TO REPRESENT AN $\ell^n$ -ISOGENY – EXAMPLE OVER $\mathbb{F}_{41^2}$

**Example:** over  $\mathbb{F}_{41^2}$ , the  $5^3$ -isogeny

$$\underbrace{(y^2 = x^3 + 1)}_E \xrightarrow{\varphi} \underbrace{(y^2 = x^3 + (16a + 10)x + (3a + 34))}_{E'} \quad \text{with } (a - 22)^2 \equiv 27 \pmod{41}$$

- the rational functions defining  $\varphi$

$$\varphi: (x, y) \mapsto \left( \frac{4x^{125} + (7a + 7)x^{124} + \dots}{10x^{124} + (38a + 38)x^{123} + \dots}, \frac{21x^{186}y + \dots}{31x^{186} + \dots} \right)$$

- (Generators of) the kernel of  $\varphi$  + Vélu's formulae

$$\ker \varphi = \{ \text{roots of } x^{62} + (6a + 6)x^{61} + \dots \} \subseteq \mathbb{F}_{41^{50}}$$

- Factor  $\varphi$  as a chain of  $\ell$ -isogenies  $\varphi_n \circ \dots \circ \varphi_1$

$$\begin{aligned} E &\xrightarrow{\varphi_1} (y^2 = x^3 + (21a + 23)x + 1) \\ &\xrightarrow{\varphi_2} (y^2 = x^3 + (34a + 11)x + (21a + 33)) \xrightarrow{\varphi_3} E' \end{aligned}$$

- The  $j$ -invariants of the codomains of  $\varphi_n, \dots, \varphi_1$

$$0 \xrightarrow{\varphi_1} 28 \xrightarrow{\varphi_2} 3 \xrightarrow{\varphi_3} 32$$

## WAYS TO REPRESENT AN $\ell^n$ -ISOGENY – EXAMPLE OVER $\mathbb{F}_{41^2}$

**Example:** over  $\mathbb{F}_{41^2}$ , the  $5^3$ -isogeny

$$\underbrace{(y^2 = x^3 + 1)}_E \xrightarrow{\varphi} \underbrace{(y^2 = x^3 + (16a + 10)x + (3a + 34))}_{E'} \quad \text{with } (a - 22)^2 \equiv 27 \pmod{41}$$

- the rational functions defining  $\varphi$

$$\varphi: (x, y) \mapsto \left( \frac{4x^{125} + (7a + 7)x^{124} + \dots}{10x^{124} + (38a + 38)x^{123} + \dots}, \frac{21x^{186}y + \dots}{31x^{186} + \dots} \right)$$

- (Generators of) the kernel of  $\varphi$  + Vélu's formulae

$$\ker \varphi = \{ \text{roots of } x^{62} + (6a + 6)x^{61} + \dots \} \subseteq \mathbb{F}_{41^{50}}$$

- Factor  $\varphi$  as a chain of  $\ell$ -isogenies  $\varphi_n \circ \dots \circ \varphi_1$

$$\begin{aligned} E &\xrightarrow{\varphi_1} (y^2 = x^3 + (21a + 23)x + 1) \\ &\xrightarrow{\varphi_2} (y^2 = x^3 + (34a + 11)x + (21a + 33)) \xrightarrow{\varphi_3} E' \end{aligned}$$

- The  $j$ -invariants of the codomains of  $\varphi_n, \dots, \varphi_1$

$$0 \xrightarrow{\varphi_1} 28 \xrightarrow{\varphi_2} 3 \xrightarrow{\varphi_3} 32$$